

Two ways to construct new sets out of old ones:

- as solution sets to systems of equations

~~~~~ limits

- by gluing sets together

~~~~~ colimits

Let I be an (index) category. An I -diagram in a category $\mathcal{C}$ is a functor $X: I \rightarrow \mathcal{C}$.

Ex 1) If I is a discrete category (i.e. only morphisms in I are the identity morphisms), then an I -diagr. $X: I \rightarrow \mathcal{C}$ determines and is determined by the set of objects $\{X_i\}_{i \in \text{ob}(I)}$.

2) If I has a single object $*$ and $\text{End}_I(*) = \text{Aut}_I(*)$ is a group G with composition

$$g \circ h = gh,$$

then an I -diagram $X: I \rightarrow \mathcal{C}$ determines and is determined by an object $X(*)$ with left G -action:

$$\left[\begin{array}{ccc} * & & X(*) \\ \downarrow h & & \downarrow h \\ gh & \xrightarrow{X} & X(*) \\ \downarrow g & & \downarrow g \\ * & & X(*) \end{array} \right] gh.$$

3) If $I = O(Y)^{op}$ is the opposite of the category of open subsets of a space Y , then an I -diagram

$$O(Y)^{op} \xrightarrow{F} \text{Sets}$$

is, by definition, a presheaf of sets on Y . If $U \subset V \subset W \subset X$, then

$$\left[\begin{array}{ccc} W & & F(W) \\ \downarrow & & \downarrow \rho_V^W \\ V & \xrightarrow{F} & F(V) \\ \downarrow & & \downarrow \rho_U^V \\ U & & F(U) \end{array} \right] \rho_U^W.$$

A limit of $X: I \rightarrow \mathcal{C}$ is a set

$$\left\{ \lim_I X \xrightarrow{\text{pr}_i} X_i \right\}_{i \in \text{ob}(I)}$$

of morphisms in $\mathcal{C}$ with the property that for every morphism $\alpha: i \rightarrow j$ in I , the diagram

$$\begin{array}{ccc} \lim_I X & \xrightarrow{\text{pr}_i} & X_i \\ & \searrow \text{pr}_j & \downarrow X(\alpha) \\ & & X_j \end{array}$$

commutes and s.t. if

$$\left\{ Y \xrightarrow{f_i} X_i \right\}_{i \in \text{ob}(I)}$$

is another set of morphism with this property, then there exists a unique morphism

$$Y \xrightarrow{f} \lim_I X_i$$

such that

$$f_i = \text{pr}_i \circ f$$

for all $i \in \text{ob}(I)$.

A colimit of $X: I \rightarrow \mathcal{C}$ is a set

$$\{ X_i \xrightarrow{\text{in}_i} \underset{I}{\text{colim}} X \}_{i \in \text{ob}(I)}$$

of morphisms in $\mathcal{C}$ with the property that for every morph. $\alpha: i \rightarrow j$ in I , the diagram

$$\begin{array}{ccc} X_i & \xrightarrow{\text{in}_i} & \underset{I}{\text{colim}} X \\ X(\alpha) \downarrow & & \nearrow \\ X_j & \xrightarrow{\text{in}_j} & \end{array}$$

commutes and s.t. if

$$\{ X_i \xrightarrow{f_i} Y \}_{i \in \text{ob}(I)}$$

is another set of morphisms with this property, then there is a unique morphism

$$\underset{I}{\text{colim}} X \xrightarrow{f} Y$$

s.t.

$$f_i = f \circ \text{in}_i$$

for all $i \in \text{ob}(I)$.

Notation: We write

$$Y \xrightarrow{(f_i)} \lim_I X$$

for the unique morphism determined by $\{f_i: Y \rightarrow X_i\}$ and

$$\operatorname{colim}_I X \xrightarrow{\Sigma f_i} Y$$

for the unique morph. determined by $\{f_i: X_i \rightarrow Y\}$.

Prop Suppose that both

$$\left\{ \lim_I X \xrightarrow{\operatorname{pr}_i} X_i \right\}_{i \in \operatorname{ob}(I)}$$

and

$$\left\{ \lim_I' X \xrightarrow{\operatorname{pr}_i'} X_i \right\}_{i \in \operatorname{ob}(I)}$$

are limits of $X: I \rightarrow \mathcal{E}$. Then

$$\lim_I X \xrightleftharpoons[(\operatorname{pr}_i)]{(\operatorname{pr}_i')} \lim_I X$$

are each others inverses. So the limit is well-defined, up to unique isomorphism. Similarly for the colimit. //

The limit of a diagram $X: I \rightarrow \text{Sets}$ of sets is given by the subset

$$\lim_I X = \prod_{i \in \text{ob}(I)} X_i$$

consisting of the tuples $(x_i)_{i \in \text{ob}(I)}$ s.t. for every morph. $\alpha: i \rightarrow j$ in I ,

$$X(\alpha)(x_i) = x_j.$$

The map $\text{pr}_i: \lim_I X \rightarrow X_i$ reads off the i th component x_i .

Dually, the colimit of $X: I \rightarrow \text{Sets}$ is given by the quotient

$$\text{colim}_I X_i = \left(\coprod_{i \in \text{ob}(I)} X_i \right) / \approx$$

where $\approx$ is the equivalence relation generated by the relation that identifies $x_i \in X_i$ and $x_j \in X_j$ if there exists a morph. $\alpha: i \rightarrow j$ in I s.t. $X(\alpha)(x_i) = x_j$. So, in general,

$$x_i \in X_i \approx x_j \in X_j$$

if there exists morphisms in I



such that the induced maps



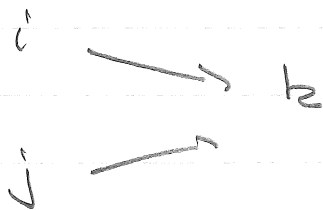
take



In general, the equivalence relation $\approx$ is very difficult to understand. It is much more well-behaved, however, if the category I is filtered, i.e.

(i) $I \neq \emptyset$;

(ii) for every pair of objects $i, j \in \text{ob}(I)$, there exists a pair of morphisms



to a common target,

(iii) for every pair of parallel morph.

$$i \begin{array}{c} \xrightarrow{\alpha} \\ \xrightarrow{\beta} \end{array} j,$$

there exists a morph. $\gamma: j \rightarrow k$
s.t. $\gamma \circ \alpha = \gamma \circ \beta: i \rightarrow k$.

Ex In examples above, (1) and (2) are not filtered except in trivial cases, but (3) is filtered. "

Lemma If $X: I \rightarrow \text{Sets}$ is a diag. of sets with I filtered, then

(i) the colimit is given by

$$\text{colim}_I X = \left(\coprod_{i \in \text{ob}(I)} X_i \right) / \sim$$

where $\sim$ is the equivalence rel. that identifies $x_i \in X_i \sim x_j \in X_j$ if there exists

$$\begin{array}{ccc} i & \xrightarrow{\alpha} & k \\ j & \xrightarrow{\beta} & k \end{array}$$

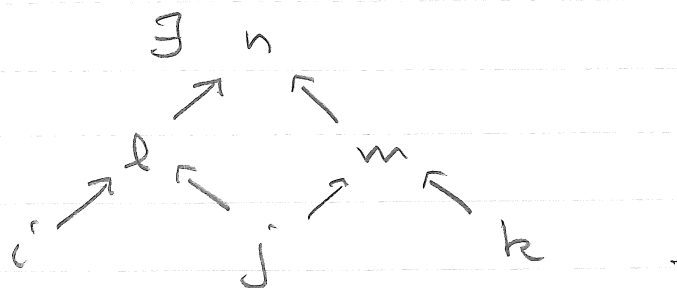
s.t. $X(\alpha)(x_i) = X(\beta)(x_j) \in X_k$.

(ii) any two elements $x_i \in X_i$ and $x_j \in X_j$ are equivalent under $\sim$

to two elements of the same X_k

(iii) two elements $x_i, x'_i \in X_i$ are equivalent under $\sim$ if and only if there exists $\alpha: i \rightarrow j$ s.t.
 $X(\alpha)(x_i) = X(\alpha)(x'_i) \in X_j$.

Pf The statement (i) is clear once we show that $\sim$ is an equivalence relation. Only transitivity is an issue, and this follows from



The statement (ii) is immediate from I being filtered. To prove (iii) suppose $x_i \sim x'_i \in X_i$. Then there exist $\alpha, \beta: i \rightarrow j$ s.t. $X(\alpha)(x_i) = X(\beta)(x'_i)$. Hence, choosing $\gamma: j \rightarrow k$ s.t.

$$\gamma \circ \alpha = \gamma \circ \beta =: \delta: i \rightarrow k,$$

we have $X(\delta)(x_i) = X(\delta)(x'_i) \in X_k$ as desired. //

Let I and J be two (index) cat-

egories, and let $X: I \times J \rightarrow \mathcal{C}$ be a diagram indexed by the product category. The universal properties of limit and colimit imply that there is a canonical morphism

$$\operatorname{colim}_I \lim_J X_{ij} \longrightarrow \lim_J \operatorname{colim}_I X_{ij}.$$

In general, it is not an isomorphism. For example, if $I = J = \emptyset$, then $\text{LHS} = \emptyset$ is an initial object in $\mathcal{C}$ while $\text{RHS} = *$ is a terminal object in $\mathcal{C}$. We say that J is finite if the sets $\text{ob}(J)$ and $\text{mor}(J)$ both are finite. The following result is extremely useful.

Prop If $X: I \times J \rightarrow \text{Sets}$ is a diagram of sets with I filtered and J finite, then the canonical map

$$\operatorname{colim}_I \lim_J X_{ij} \longrightarrow \lim_J \operatorname{colim}_I X_{ij}$$

is a bijection.

Pf Finite limits can be obtained by iterated fiber products. For

we can clearly get finite products in this way and the equalizer of $f, g: X \rightarrow Y$ is the fiber prod.

$$\begin{array}{ccc}
 X \times_{X \times Y} X & \xrightarrow{pr_1} & X \\
 \downarrow pr_2 & & \downarrow (id, f) \\
 X & \xrightarrow{(id, g)} & X \times Y
 \end{array}$$

So it suffices to show that the canonical map

$$\operatorname{colim}_I (X_i \times_{S_i} Y_i) \rightarrow (\operatorname{colim}_I X_i) \times_{(\operatorname{colim}_I S_i)} (\operatorname{colim}_I Y_i)$$

class of $(x_i, y_i) \mapsto (\text{class of } x_i, \text{class of } y_i)$

is a bijection. Construct inverse maps. Let $(\bar{x}, \bar{y})$ be an element of RHS . If $x_i \in X_i$ represents $\bar{x}$ and $y_j \in Y_j$ represents $\bar{y}$, then

$$f(x_i) \in S_i \sim g(y_j) \in S_j.$$

We pick $\alpha: i \rightarrow k$ and $\beta: j \rightarrow k$, then $x_k = X(\alpha)(x_i) \in X_k$ represents $\bar{x}$ and $y_k = Y(\beta)(y_j) \in Y_k$ represents $\bar{y}$, and

$$f(x_k) \sim g(x_k) \in S_k.$$

Finally, we pick $\gamma: k \rightarrow m$ s.t.

$$S(\gamma)(f(x_k)) = S(\gamma)(g(x_k)) \in S_m$$

and set $x_m = X(\gamma)(x_k) \in X_m$ and $y_m = Y(\gamma)(y_k) \in Y_m$. Then x_m repr. $\bar{x}$, y_m repr. $\bar{y}$, and

$$(x_m, y_m) \in X_m \times_{S_m} Y_m.$$

Now the map

$$\text{class of } (x_m, y_m) \longleftarrow (\bar{x}, \bar{y}).$$

is well-defined and inverse to the canonical map. //

Cor The underlying set of a filtered colimit of rings (resp. modules, ...) is the filtered colimit of the underlying diagram of sets.

Pf Let $A: I \rightarrow \text{Rings}$ be a filtered diagram of rings. The colimit of the underlying diagram of sets inherits a ring structure:

$$(\operatorname{colim}_I A_i) \times (\operatorname{colim}_I A_i)$$

$$\xleftarrow{\sim} \operatorname{colim}_I (A_i \times A_i)$$

$$\xrightarrow{+,\cdot} \operatorname{colim}_I A_i$$

Therefore, this is the colimit in the category of rings. "

Ex This is not true if I is not filtered. For example, if $I = \emptyset$, then the colimit is the initial object, which in Sets is $\emptyset$, but in Rings is $\mathbb{Z}$. "

Let X be a space, let $X^\wedge$ be the category of presheaves of sets on X , and let $X^\sim \subset X^\wedge$ be the full subcategory of sheaves.

Thm The canonical inclusion functor

$$X^\sim \longleftrightarrow X^\wedge$$

admits a left adjoint

$$X^\wedge \xrightarrow{a} X^\sim$$

and a preserves finite limits.

The functor a is called sheafification.
If $\{U_\alpha \rightarrow U\}_{\alpha \in A}$ is a covering, define

$$H^0(\{U_\alpha \rightarrow U\}, F) \rightarrow \prod_{\alpha \in A} F(U_\alpha) \xrightarrow[\substack{\downarrow \\ \alpha, \alpha' \in A}]{\substack{\downarrow \\ \alpha, \alpha' \in A}} \prod_{\alpha, \alpha' \in A} F(U_\alpha \cap U_{\alpha'})$$

to be an equalizer. We define a map of coverings

$$\{U_\alpha \rightarrow U\}_{\alpha \in A} \rightarrow \{V_\beta \rightarrow U\}_{\beta \in B}$$

to be a map $f: A \rightarrow B$ s.t.

$$\begin{array}{ccc} U_\alpha & \hookrightarrow & V_{f(\alpha)} \\ \downarrow & & \downarrow \\ U & = & U \end{array}$$

It induces

$$H^0(\{V_\beta \rightarrow U\}, F) \xrightarrow{f^*} H^0(\{U_\alpha \rightarrow U\}, F).$$

Lemma If $f, g: \{U_\alpha \rightarrow U\} \rightarrow \{V_\beta \rightarrow U\}$ are any two maps, then $f^* = g^*$.

Pf Exercise. //

For $U \subset X$ open, let $I(U)$ be the category with objects coverings $\{U_\alpha \rightarrow U\}$ of U and with a single

morphism $\{U_\alpha \rightarrow U\} \rightarrow \{V_\beta \rightarrow U\}$ is there exists a map of coverings from $\{U_\alpha \rightarrow U\}$ to $\{V_\beta \rightarrow U\}$. The lemma shows that

$$I(U)^{op} \xrightarrow{H^0(-, F)} \text{Sets}$$

is a well-defined functor, and we define

$$F^+(U) = \operatorname{colim}_{I(U)^{op}} H^0(-, F)$$

It is a filtered colimit, since

$$\begin{array}{ccc} \{U_\alpha \rightarrow U\}_\alpha & \leftarrow & \\ & \nwarrow & \\ \{V_\beta \rightarrow U\}_\beta & \leftarrow & \{U_\alpha \cap V_\beta \rightarrow U\}_{\alpha, \beta} \end{array}$$

and since there are no parallel morphisms. Moreover, if $U \rightarrow V$, then

$$I(U) \longleftarrow I(V)$$

$$\{U \cap V_\alpha \rightarrow U\} \longleftarrow \{V_\alpha \rightarrow V\}$$

which, in turn, induces

$$F^+(U) \longleftarrow F^+(V).$$

So F^+ is a presheaf on X . We say that a presheaf on X is separated if for every covering $\{U_\alpha \rightarrow U\}_{\alpha \in A}$,

$$F(U) \xrightarrow{i} \prod_{\alpha \in A} F(U_\alpha)$$

is injective. The proofs of the following lemmas are given in Artin's "Grothendieck Topologies."

Lemma For every presheaf F on X , the presheaf F^+ is separated.

Lemma For every separated presheaf F on X , F^+ is a sheaf.

Proof of Theorem: There is a canonical natural map of presheaves $\eta: F \rightarrow F^+$ which is an isomorphism, if F is separated. Hence, if $f: F \rightarrow G$ is a map of presheaves, the diagram

$$\begin{array}{ccccc} F & \xrightarrow{\eta} & F^+ & \xrightarrow{\eta} & F^{++} \\ \downarrow f & & \downarrow f^+ & & \downarrow f^{++} \\ G & \xrightarrow[\sim]{\eta} & G^+ & \xrightarrow[\sim]{\eta} & G^{++} \end{array}$$

shows that $a(F) = F^{++}$ is left

adjoint to $\hat{c}$. By construction, $\hat{c}$ preserves finite limits. //

Prop (Kan extension) let $\mathcal{C}$ be a category which admits all limits and colimits, and let $u: I \rightarrow J$ be a functor. Then the functor

$$\mathcal{C}^J \xrightarrow{u^*} \mathcal{C}^I$$

$$F \longmapsto F \circ u$$

has both a left adjoint $u_!$ and a right adjoint u_* given by

$$(u_! F)(j) = \operatorname{colim}_{(i, u(i) \rightarrow j)} F(i)$$

$$(u_* F)(j) = \lim_{(i, j \rightarrow u(i))} F(i)$$

Pf Exercise. //

let X be a space and let $O(X)$ be the category of open subsets:

$$\operatorname{ob} O(X) = \{ U \subset X \text{ open} \}$$

$\operatorname{mor} O(X)$: a single morph. $U \rightarrow V$ if $U \subset V$.

A continuous map $f: X \rightarrow Y$ induces the functor

$$O(X) \xleftarrow{f^{-1}} O(Y).$$

Hence we have functors

$$\begin{array}{ccc} X^\wedge & \begin{array}{c} \xleftarrow{(f^{-1})!} \\ \xrightarrow{(f^{-1})^*} \\ \xleftarrow{(f^{-1})_*} \end{array} & Y^\wedge \end{array}$$

given by

$$\begin{aligned} (f_p G)(V) &:= (f^{-1})^*(G)(V) \\ &= G(f^{-1}(V)) \end{aligned}$$

$$\begin{aligned} (f^p F)(U) &:= (f^{-1})_!(F)(U) \\ &= \operatorname{colim}_{(V, U \subset f^{-1}(V))} F(V) \end{aligned}$$

The direct image functor f_p preserves sheaves in that

$$\begin{array}{ccc} X^\wedge & \xrightarrow{f_p} & Y^\wedge \\ \uparrow i_X & & \uparrow i_Y \\ X^\sim & \xrightarrow{f_*} & Y^\sim \end{array}$$

commutes for a unique functor f_* .
The functor f^P , in general, does not preserve sheaves, so we define the inverse image functor

$$X^\sim \xleftarrow{f^*} Y^\sim$$

to be the composite $a_x f^P i_y$.

Cor The functors (f^*, f_*) form an adjoint pair, and f^* preserves finite limits.

Pf Prove adjointness:

$$\begin{aligned} \text{Hom}_{X^\sim}(f^* F, G) &= \text{Hom}_{X^\sim}(a_x f^P i_y F, G) \\ &\xrightarrow{\sim} \text{Hom}_{X^\wedge}(f^P i_y F, i_x G) \\ &\xrightarrow{\sim} \text{Hom}_{Y^\wedge}(i_y F, f_P i_x G) \\ &= \text{Hom}_{Y^\wedge}(i_y F, i_y f_* G) \\ &= \text{Hom}_{Y^\sim}(F, f_* G). \end{aligned}$$

The colimit defining f^P is filtered, so f^P pres. finite limits. So does a_x and i_y preserves all limits.
Hence, f^* pres. finite limits. //